# Software project 'miraculix': <br> Efficient computations with large genomic datasets 

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## Idea and motivation behind 'miraculix'

- miraculix is a library, not a standalone package
- goals
- improving time (and/or memory) critical parts of code used in genetics
- detection of hardware during run-time
- close cooperations with partners at Wageningen and Göttingen
- code published irregularly on github
- CPU solutions are represented here GPU solutions will mainly be presented in the next talk


## Integration of miraculix

MoBPS uses miraculix for

- Compressed storage of genomic datasets (twobit format)
- Efficient calculation of the genomic relationship matrix
- Computation of BVE for gBLUP models (Cholesky decomposition)
$\rightarrow$ Breeding program simulation is si-
gnificantly accelerated
(Pook et al., 2021)

MiXBLUP uses miraculix for

- Acceleration of genotype matrix multiplications
- Faster iteration times in solving single-step models
$\rightarrow \mathrm{BVE}$ in single-step models is substantially faster
(Freudenberg et al., 2023b)


## Mixed Model Equations

- Goal: Solving single-step Mixed Model Equations, e.g., ssGBLUP:

$$
\left(\begin{array}{cc}
X^{\top} R^{-1} X & X^{\top} R^{-1} W \\
W^{\top} R^{-1} X & W^{\top} R^{-1} W+H^{-1}
\end{array}\right)\binom{\hat{b}}{\hat{u}}=\binom{X^{\top} R^{-1} y}{W^{\top} R^{-1} y}
$$

- Solver software uses iterative algorithms (e.g., PCG)
- Each iteration requires multiplication of the coefficient matrix
- Coefficient matrix involves the genotype matrix $Z$ and its transposed $Z^{T}$
- for ease, $Z[\ldots] Z^{\top} V$, but also $Z Z^{\top}$ will be considered in the following


## Time Results $Z Z^{\top}$

Time measurements for GRM calculation with 10k individuals


Hardware: Xeon Platinum 8368
Time for reading and writing is not measured when miraculix is used. $\rightarrow$ An additive constant in the computing times of miraculix is missing.

## Computing times for $Z\left(Z^{\top} V\right)$



Left:

- 50k individuals
- small/medium/large = 102k/700k/3100k snps
- AMD Milan EPYC 7513 (20 cores)

SNP matrix $Z$ with indiv=50k calculating $\mathbf{Z}\left(\mathbf{Z}^{\wedge} \mathrm{t} \mathrm{V}\right)$ with AVX2


Right:

- 50k individuals
- small/medium/large $=$ 102k/700k/3100k snps
- XEON 6230


## Frame conditions

- libraries for double-double scalar products are well-developed
- current approaches decompress packed 2-bit representation (of plink)
- transformation of data is always worth, as soon as computing time is of smaller order.
- additional memory are considered acceptable (once or twice the original size)


## Hash table calculating $Z Z^{\top}$ (based on plink coding)

## Basis: SIMD command mm_shuffle_epi8

- hash table: 16 entries, 1 byte each
- addressed simultaneously by lower $\frac{1}{2}$ byte of each byte in the register
- 16 hash table look-ups at once for SIMD (64 for AVX512)

Goal: find operator o with identifiable results

| * |  | 0 | 1 | 2 | $\stackrel{\text { hash table }}{\leftarrow}$ | $\bigcirc$ |  | 000 | $\begin{gathered} 1 \\ 10_{p} \end{gathered}$ | $\begin{gathered} 2 \\ 11_{p} \end{gathered}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $00_{p}$ | $10_{p}$ | $11_{p}$ |  |  |  |  |  |  |
| 0 | $00_{p}$ | 0 | 0 | 0 |  | 0 | $00_{p}$ | 00 | 00 | 00 |
| 1 | $10_{p}$ | 0 | 1 | 2 |  | 1 | $10_{p}$ | 00 | 10 | 01 |
| 2 | $11_{p}$ | 0 | 2 | 4 |  | 2 | $11_{p}$ | 00 | 01 | 11 |

operator $\circ$ is a composition of bitwise \&, |, >>, and a substraction

Idea for calculating $Z\left(Z^{\top} V\right)$ for AVX2 only
Basis: Hash table in the L1 cache

- hash table: 243 entries, 1 double each

Mathematical background
$Z_{1}, \ldots, Z_{5} \in\{0,1,2\}$ : arbitrary SNP values
$V_{1}, \ldots, V_{5} \in \mathbb{R}$, fixed
Scalar product of $\left(Z_{1}, \ldots, Z_{5}\right)$ with $\left(V_{1}, \ldots, V_{5}\right)$,

$$
Z_{1} V_{1}+\ldots+Z_{5} V_{5}
$$

takes only $3^{5}=243$ different values.

Idea

- compress 5 SNP values into 1 byte
- this byte addresses the result of the scalar product in the hash table


## References

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